

Static and Dynamic Analysis

17-313, Foundations of Software Engineering, Spring 2023



Learning Goals

- Gain an understanding of the relative strengths and weaknesses of static and dynamic analysis
- Examine several popular analysis tools and understand their use cases
- Understand how analysis tools are used in large open source software



Activity: Analyze the Python program statically

```
def n2s(n: int, b: int):
 if n \le 0: return '0'
 r = "
 while n > 0:
  u = n \% b
  if u >= 10:
    u = chr(ord('A') + u-10)
  n = n // b
  r = str(u) + r
 return r
```

- 1. What are the set of data types taken by variable `u` at any point in the program?
- 2. Can the variable `u` be a negative number?
- 3. Will this function always return a value?
- 4. Can there ever be a division by zero?
- 5. Will the returned value ever contain a minus sign '-'?



What static analysis can and cannot do

- Type-checking is well established
 - Set of data types taken by variables at any point
 - Can be used to prevent type errors (e.g. Java) or warn about potential type errors (e.g. Python)
- Checking for problematic patterns in syntax is easy and fast
 - o Is there a comparison of two Java strings using `==`?
 - o Is there an array access `a[i]` without an enclosing bounds check for `i`?
- Reasoning about termination is impossible in general
 - Halting problem
- Reasoning about exact values is hard, but conservative analysis via abstraction is possible
 - Is the bounds check before `a[i]` guaranteeing that `I` is within bounds?
 - Can the divisor ever take on a zero value?
 - Could the result of a function call be `42`?
 - Will this multi-threaded program give me a deterministic result?
 - Be prepared for "MAYBE"
- Verifying some advanced properties is possible but expensive
 - CI-based static analysis usually over-approximates conservatively



The Bad News: Rice's Theorem

Every static analysis is necessarily incomplete, unsound, undecidable, or a combination thereof

"Any nontrivial property about the language recognized by a Turing machine is undecidable."

Henry Gordon Rice, 1953

Static Analysis is well suited to detecting certain defects

- Security: Buffer overruns, improperly validated input...
- Memory safety: Null dereference, uninitialized data...
- Resource leaks: Memory, OS resources...
- API Protocols: Device drivers; real time libraries; GUI frameworks
- Exceptions: Arithmetic/library/user-defined
- Encapsulation:
 - Accessing internal data, calling private functions...
- Data races:
 - Two threads access the same data without synchronization

Activity: Analyze the Python program dynamically

```
def n2s(n: int, b: int):
 if n \le 0: return '0'
 r = "
 while n > 0:
  u = n \% b
  if u >= 10:
    u = chr(ord('A') + u-10)
  n = n // b
  r = str(u) + r
 return r
```

- 1. What are the set of data types taken by variable `u` at any point in the program?
- 2. Did the variable `u` ever contain a negative number?
- 3. For how many iterations did the while loop execute?
- 4. Was there ever be a division by zero?
- 5. Did the returned value ever contain a minus sign '-'?





Dynamic analysis reasons about program executions

- Tells you properties of the program that were definitely observed
 - Code coverage
 - Performance profiling
 - Type profiling
 - Testing
- In practice, implemented by program instrumentation
 - Think "Automated logging"
 - Slows down execution speed by a small amount



Static Analysis

- Requires only source code
- Conservatively reasons about all possible inputs and program paths
- Reported warnings may contain false positives
- Can report all warnings of a particular class of problems
- Advanced techniques like verification can prove certain complex properties, but rarely run in CI due to cost

Dynamic Analysis

- Requires successful build + test inputs
- Observes individual executions
- Reported problems are real, as observed by a witness input
- Can only report problems that are seen. Highly dependent on test inputs. Subject to false negatives
- Advanced techniques like symbolic execution can prove certain complex properties, but rarely run in CI due to cost



Static Analysis Tools

Tools for Static Analysis



























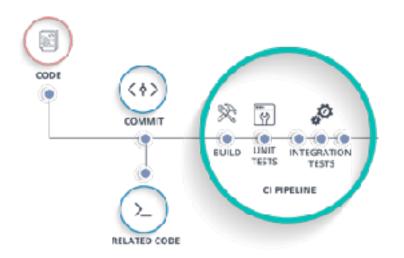
Static analysis can be applied to all attributes

- Find bugs
- Refactor code
- Keep your code stylish!
- Identify code smells
- Measure quality
- Find usability and accessibility issues
- Identify bottlenecks and improve performance





Static analysis is a key part of continuous integration









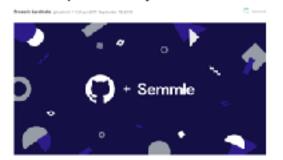






Static analysis is a growing industry

GitHub acquires code analysis tool Semmle









Snyk Secures \$150M, Snags \$1B Valuation

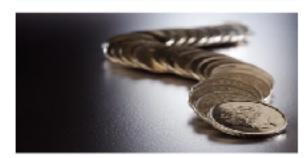












Snyl, a developer focused secure, startup that and identifies vulnerabilities in open source applications. assounced a \$150 million Series O funding round today. This brings the company's total investment to \$250 million alongside reports that out the company's valuation at more than \$1 billion.





Static analysis is also integrated into IDEs









```
Cppcorequiteines.cpp
       I// To enable only C++ Core Guidelines checks
       // go to Settings/Preferences | Editor | Inspections | C/C++ | Clang-Tidy
       // and provide: --, oppcoreguidelines-+ is options
        weid fill_pointer(int+ arr, const int num) {
            for(int 1 = 0: 1 < num; ++1) (
                arrial = *:
     Do not use pointer arithmetic
        weid fill arroy(int ind) (
            ist arr[3] - (1,3,3):
13
            arr[ind] = 0:
14
15
16
        weid cast away const(canst int& magic num)
            const_east<imt&>(magic_num) = 42;
19
20
```

```
ni Example.m.
     3. T. C. ** B #
2. Object allocated on line 13 is no longer referenced after this point and has a ... : (4 | +) (Done)
 10
      void-forbiat s, int vl {
         id obj - [[MSString alloc] init];
                                 Witholl of the an Olde Twe Colded with a of the Common Livering reference.
           case 🥸:
 16
             [obj release];
 LT.
             break;
           case 1:
 130
 19
                       febj autoreleasel:
 70
            break;
21
           default:
             broak;
            Chieco allocated on line 13 taing long an reference diafter this point and has a recain count of +1 toblets legislati.
```

What makes a good static analysis tool?

- Static analysis should be fast
 - Don't hold up development velocity
 - This becomes more important as code scales
- Static analysis should report few false positives
 - Otherwise developers will start to ignore warnings and alerts, and quality will decline
- Static analysis should be continuous
 - Should be part of your continuous integration pipeline
 - Diff-based analysis is even better -- don't analyse the entire codebase; just the changes
- Static analysis should be informative
 - Messages that help the developer to quickly locate and address the issue
 - Ideally, it should suggest or automatically apply fixes



Linters

Cheap, fast, and lightweight static source analysis





Linters for Maintainability

Use linters to improve maintainability Why? We spend more time reading code than writing it.

- Developers spend most of their time maintaining code
 - Various estimates of the exact %, some as high as 80%
- Code is ownership is usually shared
- The original owner of some code may move on
- Code conventions make it easier for other developers to quickly understand your code



Use Style Guidelines to facilitate communication

- Indentation
- Comments
- Line length
- Naming
- Directory structure
- ...







Guidelines are inherently opinionated, but **consistency** is the important point. Agree to a set of conventions and stick to them.



Use linters to enforce style guidelines

Don't rely on manual inspection during code review!













Example: CheckStyle



```
cmodule name="WhitespaceArcund">
 cproperty name="allowEmptyConstructors" value="true'/>
 <procerty name="allowEmptyLambdas" value="true"/>
 cproperty name="allowEmptyMethods" value="true"/>
 cproperty name="allowEmptyTypes" value="true"/>
 cprocerty name="allowEmptyLoops" value="true"/>
 cprocerty name="ignoreEnhancedForColon" value="false"/>
 cproperty name="tokens"
          EXOR ASSIGN, COLON, CIV. DIV ASSIGN, CO WHILE, EQUAL, CE, CT, LAMBDA, LAND.
              LCURLY, LE, LITERAL CATCH, LITERAL DO, LITERAL ELSE, LITERAL FINALLY,
              LITERAL FOR, LITERAL IF, LITERAL RETURN, LITERAL SWITCH, LITERAL SYNCHRONIZED,
              LITERAL_TRY, LITERAL_WHILE, LOR, LT, MINUS, MINUS_ASSIGN, MOD, MOD_ASSIGN,
              NOT_EQUAL, PLUS_ASSIGN, QUESTION, RCURLY, SL, SLIST, SL_ASSIGN, SR,
              SR_ASSIGN, STAR, STAR_ASSIGN, LITERAL_ASSERT, TYPE_EXTENSION_AND*/>
 <message key="ws.notFellowed"</pre>
         value="WhitespaceAround: ''{0}'' is not followed by whitespace. Empty blocks may only
 <message key="ws.notPreceded"</pre>
         value="WhitespaceAround: ''{0}'' is not preceded with whitespace."/>
</medule>
```

...



```
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```

```
public boolean equals(Object o) {
   if (o == this)
        return true;

private Board cur_board;
private Board blank_board;

public static void main(String[] args) throws Exception {
    TicTacToe tictactoe = new TicTacToe();
}
```

Integrate style checking into your CI

```
plugins {
   id 'checkstyle'
}
...
checkstyle {
   ignoreFailures = true
   toolVersion = "6.7"
   sourceSets = [sourceSets.main]
}
```







```
name: Tox lint checking
on: [pull_request]
jobs:
 build:
   runs-on: ubuntu-28.34
   steps:
   - uses: actions/checkoutev2
   - name: Install Python
     uses: actions/setup-python@v2
     with:
       sython-version: 3.9.5
     name: Install pipeny
     run; pip install pipenv--2021.5.29
   - id: (ache-piperv
     uses: actions/cache@v2
     with:
        path: ~/.local/share/virtualenvs
        key: ${{ runrer.os }}-pipenv-${{ hashFiles('**/Pipfile.lock') }}
   - name: Install package
     if: steps.cache-pipenv.outputs.cache-bit != 'true'
     run:
        sigeny install -- dev
   - nane; Flake0
     run: piperv run flakeð src
    - nane: MyPy
     run: piperv run mypy src
```



Automatically reformat your existing code Developer time is valuable!







Take Home Message: Style is an easy way to improve readability

- Everyone has their own opinion (e.g., tabs vs. spaces)
- Agree to a convention and stick to it
 - Use continuous integration to enforce it
- Use automated tools to fix issues in existing code





Pattern-Based Static Analyzers

Cheap and fast tools that scan Abstract Syntax Trees for common developer mistakes known as patterns













SpotBugs

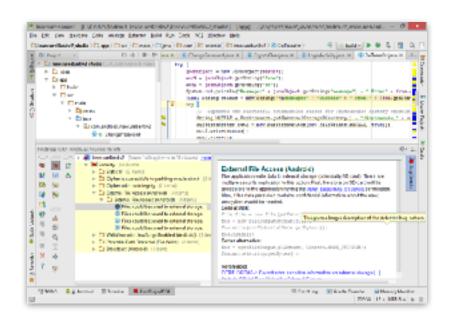
- Bad Practice
- Correctness
- Performance
- Internationalization
- Malicious Code
- Multithreaded Correctness
- Security
- Dodgy Code

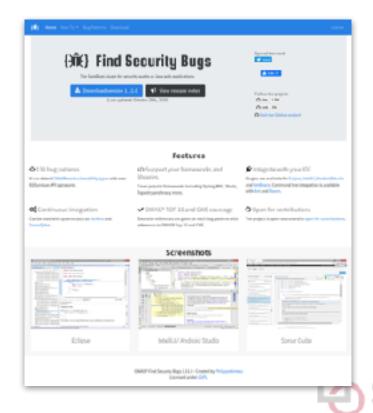






SpotBugs can be extended with plugins





Bad Practice:

```
String x = new String("Foo");
String y = new String("Foo");

if (x == y) {
    System.out.println("x and y are the same!");
} else {
    System.out.println("x and y are different!");
}
```



Bad Practice: ES_COMPARING_STRINGS_WITH_EQ Comparing strings with ==

```
String x = new String("Foo");
String y = new String("Foo");

if (x == y) {
   if (x.equals(y)) {
      System.out.println("x and y are the same!");
   } else {
      System.out.println("x and y are different!");
}
```

Performance:

```
public static String repeat(String string, int times)
{
   String output = string;
   for (int i = 1; i < times; ++i) {
      output = output + string;
   }
   return output;
}</pre>
```

Performance: SBSC_USE_STRINGBUFFER_CONCATENATION Method concatenates strings using + in a loop

```
public static String repeat(String string, int times)
{
   String output = string;
   for (int i = 1; i < times; ++i) {
      output = output + string;
   }
   return output;
}</pre>
```

The method seems to be building a String using concatenation in a loop. In each iteration, the String is converted to a StringBuffer/StringBuilder, appended to, and converted back to a String. **This can lead to a cost quadratic in the number of iterations, as the growing string is recopied in each iteration.**

Performance: SBSC_USE_STRINGBUFFER_CONCATENATION Method concatenates strings using + in a loop

```
public static String repeat(String string, int times)
{
   StringBuffer output = new StringBuffer(string);
   for (int i = 1; i < times; ++i) {
      output.append(string);
   }
   return output.toString();
}</pre>
```

Performance: SBSC_USE_STRINGBUFFER_CONCATENATION Method concatenates strings using + in a loop

```
public static String repeat(String string, int times)
{
  int length = string.length() * times;
  StringBuffer output = new StringBuffer(length);
  for (int i = 0; i < times; ++i) {
    output.append(string);
  }
  return output.toString();
}</pre>
```

Correctness: Lots of issues here!

```
public class QwicsXid implements Xid {
 private byte[] globalTransactionId;
 private byte[] branchQualifier;
 private int formatld;
 @Override
 public byte[] getBranchQualifier() {
  return this.branchQualifier;
 @Override
 public int getFormatId() {
  return this.getFormatId();
 @Override
 public byte[] getGlobalTransactionId() {
  return this.getGlobalTransactionId();
```

Description	Renounce	Path
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Correctness:

```
@Override
public Connection getConnection() throws SQLException {
 QwicsConnection con = new QwicsConnection(host, port);
 try {
  con.open();
 } catch (Exception e) {
  new SQLException(e);
 return con;
```

Correctness:

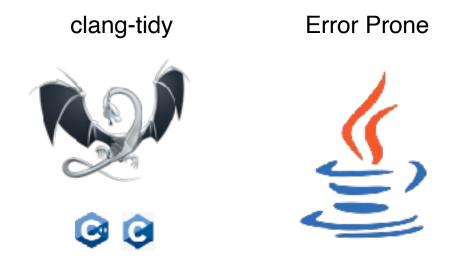
```
@Override
public Connection getConnection() throws SQLException {
 QwicsConnection con = new QwicsConnection(host, port);
 try {
  con.open();
 } catch (Exception e) {
  throw new SQLException(e);
 return con;
```

What are some of the problems with SpotBugs?



Google: Move static checks to the compiler

Developers can ignore warnings, but they can't ignore build errors





New languages have embraced the same idea Code smells will cause the build to fail (e.g., dead code)









Challenges

- The analysis must produce zero false positives
 - Otherwise developers won't be able to build the code!
- The analysis needs to be really fast
 - o Ideally < 100 ms
 - If it takes longer, developers will become irritated and lose productivity
- You can't just "turn on" a particular check
 - Every instance where that check fails will prevent existing code from building
 - There could be thousands of violations for a single check across large codebases



Challenges

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Solution: Automatically patch existing bugs

```
public class StringIsEmpty {
 @BeforeTemplate
 boolean equalsEmptyString(String string) {
  return string.equals("");
 @BeforeTemplate
 boolean lengthEquals0(String string) {
  return string.length() == 0;
 @AfterTemplate
 @AlsoNegation
 boolean optimizedMethod(String string) {
  return string.isEmpty();
```

@BeforeTemplate finds String expressions that match the body of the method.

@AfterTemplate rewrites matching String expressions to match the body of the method.

Solution: Automatically patch existing bugs

```
boolean b = someChained().methodCall().returning. AString().length() == 0;

boolean b = someChained().methodCall().returning. AString().isEmpty();
```



Summary: Linters

- Linters are cheap and fast static analysis tools!
- Style checkers can improve readability of code
- Pattern-based bug detectors catch common developer mistakes
 - Code smells, performance issues, correctness, ...
 - They don't know the intent of the program, leading to occasional false positives
 - They reveal issues that are genuine, but which we don't sufficiently care about
 - The best tools automatically fix detected issues
 - Each developer mistake needs its own analyzer / AST checker
 - They complement but don't replace testing



Java Checker Framework Uses annotations to detect common errors

- Uses a conservative analysis to prove the absence of certain defects *
 - Null pointer errors, uninitialized fields, certain liveness issues, information leaks, SQL injections, bad regular expressions, incorrect physical units, bad format strings, ...
 - C.f. SpotBugs which makes no safety guarantees
 - Assuming that code is annotated and those annotations are correct
- Uses annotations to enhance Java's type system



Annotations can be applied to types and declarations

```
// return value
@InternedString intern() { ... }
// parameter
int compareTo(@NonNullString other) { ... }
// receiver ("this" parameter)
String to String(@TaintedMyClass this) { ... }
// generics: non-null list of interned Strings
@NonNullList<@InternedString> messages;
// arrays: non-null array of interned Strings
@InternedString @NonNull[] messages:
// cast
myDate = (@InitializedDate) beingConstructed;
```

Detecting null pointer exceptions

- @Nullable indicates that an expression may be null
- @NonNull indicates that an expression must never be null
 - Rarely used because @NonNull is assumed by default
 - See documentation for other nullness annotations
- Guarantees that expressions annotated with @NonNull will never evaluate to null, forbids other expressions from being dereferenced



```
import org.checkerframework.checker.nullness.qual.*;
public class NullnessExampleWithWarnings {
    public void example() {
      @NonNull String foo = "foo";
      String bar = null;
      foo = bar;
    }
}
```

```
import org.checkerframework.checker.nullness.qual.*;

public class NullnessExampleWithWarnings {
    public void example() {
        @NonNull String foo = "foo";
        String bar = null;

    foo = bar;
    }
}
```

```
import org.checkerframework.checker.nullness.qual.*;
public class NullnessExampleWithWarnings {
 public void example() {
  @NonNull String foo = "foo";
  String bar = null; // @Nullable
                                           bar is refined to @NonNull
  if (bar != null) {
   foo = bar;
```

Is there a bug?

```
public String getDay(int dayIndex) {
 String day = null;
 switch (dayIndex) {
  case 0: day = "Monday";
  case 1: day = "Tuesday";
  case 2: day = "Wednesday";
  case 3: day = "Thursday";
 return day;
public void example() {
 @NonNull String dayName = getDay(4);
 System.out.println("Today is " + dayName);
```

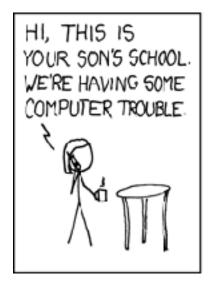
Is there a bug? Yes.

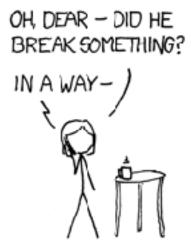
```
public String getDay(int dayIndex) {
 String day = null;
 switch (dayIndex) {
  case 0: day = "Monday";
  case 1: day = "Tuesday";
  case 2: day = "Wednesday";
  case 3: day = "Thursday";
 return day;
                                   Error: [return.type.incompatible] incompatible types in return.
                                     type of expression: @Initialized @Nullable String
                                     method return type: @Initialized @NonNull String
public void example() {
 @NonNull String dayName = getDay(4);
 System.out.println("Today is " + dayName);
```

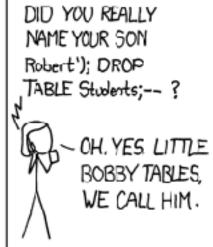
Taint Analysis

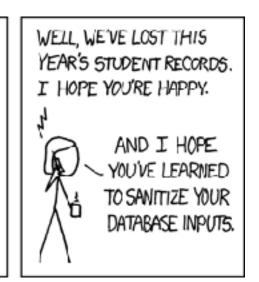
Prevents untrusted (tainted) data from reaching sensitive locations (sinks)

- Tracks flow of sensitive information through the program
- Tainted inputs come from arbitrary, possibly malicious sources
 - User inputs, unvalidated data
- Using tainted inputs may have dangerous consequences
 - Program crash, data corruption, leak private data, etc.
- We need to check that inputs are sanitized before reaching sensitive locations









```
void processRequest() {
   String input = getUserInput();
   String query = "SELECT ... " + input;
   executeQuery(query);
}
```

Tainted input arrives from an untrusted source

```
void processRequest() {
   String input = getUserInput();
   String query = "SELECT ... " + input;
   executeQuery(query);
}
```

Tainted input flows to a sensitive sink

Taint is removed by sanitizing the data

```
void processRequest() {
   String input = getUserInput();
   input = sanitizeInput(input);
   String query = "SELECT ... " + input;
   executeQuery(query);
}
```

We can now safely execute query on untainted data

Taint Checker: @Tainted and @Untainted

```
void processRequest() {
 @Tainted String input = getUserInput();
 executeQuery(input);
public void executeQuery(@Untainted String input) {
 // ...
@Untainted public String validate(String userInput) {
 // ...
```

Taint Checker: @Tainted and @Untainted

```
Indicates that data is tainted
void processRequest() {
 @Tainted String input = getUserInput();
 executeQuery(input);
                                                        Argument must be untainted
public void executeQuery(@Untainted String input) {
 // ...
                                            Guarantees that return value is untainted
@Untainted public String validate(String userInput) {
 // ...
```

Taint Checker: @Tainted and @Untainted

```
Indicates that data is tainted
void processRequest() {
 @Tainted String input = getUserInput();
 executeQuery(input);
                                                       Argument must be untainted
public void executeQuery(@Untainted String input) {
 // ...
                                           Guarantees that return value is untainted
@Untainted public String validate(String userInput) {
                               Does this compile?
```

```
void processRequest() {
 @Tainted String input = getUserInput();
 input = validate(input);
 executeQuery(input);
                                                Input becomes @Untainted
public void executeQuery(@Untainted String input) {
 // ...
@Untainted public String validate(String userInput) {
 //
```

Does this program compile?

```
void processRequest() {
    @Tainted String input = getUserInput();
    if (input.equals("little bobby drop tables")) {
        input = validate(input);
    }
    executeQuery(input);
}
```

Does this program compile? No.

```
void processRequest() {
    @Tainted String input = getUserInput();
    if (input.equals("little bobby drop tables")) {
        input = validate(input); // @Untainted
    }
    executeQuery(input); // @Tainted
}
```



Remember the Mars Climate Orbiter incident from 1999?



NASA's Mars Climate Orbiter (cost of \$327 million) was lost because of a discrepancy between use of metric unit Newtons and imperial measure Poundforce.

Units Checker identifies physical unit inconsistencies

- Guarantees that operations are performed on the same kinds and units
- Kind annotations
 - @Acceleration, @Angle, @Area, @Current, @Length, @Luminance,
 @Mass, @Speed, @Substance, @Temperature, @Time
- SI unit annotation
 - @m, @km, @mm, @kg, @mPERs, @mPERs2, @radians, @degrees,
 @A, ...



```
import static org.checkerframework.checker.units.UnitsTools.m;
import static org.checkerframework.checker.units.UnitsTools.mPERs;
import static org.checkerframework.checker.units.UnitsTools.s;
void demo() {
 @m int x;
x = 5 * m;
 @m int meters = 5 * m;
 @s int seconds = 2 * s;
 @mPERs int speed = meters / seconds;
 @m int foo = meters + seconds;
 @s int bar = seconds - meters;
```

```
import static org.checkerframework.checker.units.UnitsTools.m;
import static org.checkerframework.checker.units.UnitsTools.mPERs;
import static org.checkerframework.checker.units.UnitsTools.s;
                                      @m indicates that x represents meters
void demo() {
 @m int x;
x = 5 * m;
                                                    To assign a unit, multiply appropriate
                                                    unit constant from UnitTools
 @m int meters = 5 * m;
 @s int seconds = 2 * s;
 @mPERs int speed = meters / seconds;
 @m int foo = meters + seconds:
 @s int bar = seconds - meters;
```

Does this program compile?

```
import static org.checkerframework.checker.units.UnitsTools.m;
import static org.checkerframework.checker.units.UnitsTools.mPERs;
import static org.checkerframework.checker.units.UnitsTools.s;
                                      @m indicates that x represents meters
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 @s int seconds = 2 * s;
 @mPERs int speed = meters / seconds;
 @m int foo = meters + seconds:
 @s int bar = seconds - meters;
```

Does this program compile? No.

```
import static org.checkerframework.checker.units.UnitsTools.m;
import static org.checkerframework.checker.units.UnitsTools.mPERs;
import static org.checkerframework.checker.units.UnitsTools.s;
```

```
void demo() {
  @m int x;
  x = 5 * m;

  @m int meters = 5 * m;
  @s int seconds = 2 * s;

  @mPERs int speed = meters / seconds;
  @m int foo = meters + seconds;
  @s int bar = seconds - meters;
}
```

Addition and subtraction between meters and seconds is physically meaningless

Checker Framework: Limitations

- Can only analyze code that is annotated
 - Requires that dependent libraries are also annotated
 - o Can be tricky, but not impossible, to retrofit annotations into existing codebases
- Only considers the signature and annotations of methods
 - Doesn't look at the implementation of methods that are being called
- Dynamically generated code
 - Spring Framework
- Can produce false positives!
 - Byproduct of necessary approximations

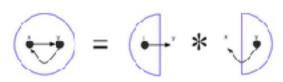


Infer: What if we didn't need annotations?

- Focused on memory safety bugs
 - Null pointer dereferences, memory leaks, resource leaks, ...
- Compositional interprocedural reasoning
 - Based on separation logic and bi-abduction
- Scalable and fast
 - Can run incremental analysis on changed code
- Does not require annotations
- Supports multiple languages
 - Java, C, C++, Objective-C
 - Programs are compiled to an intermediate representation



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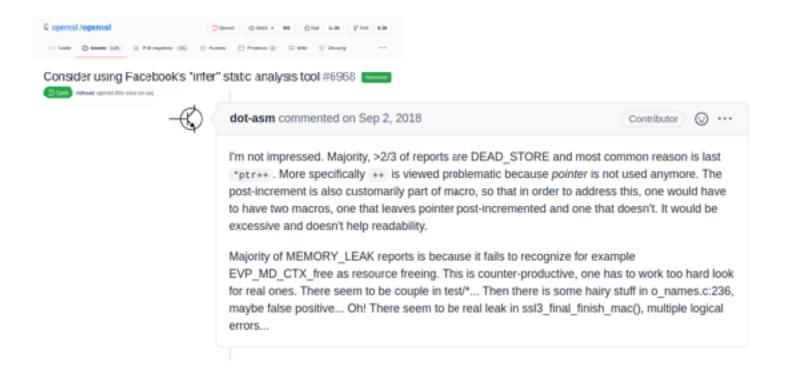


Infer: Hello World!

```
// Hello.java
class Hello {
  int test() {
    String s = null;
    return s.length();
  }
}
```

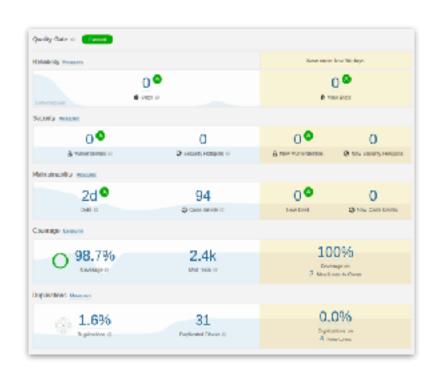
```
$ infer run -- javac Hello.java
...
Hello.java:5: error: NULL_DEREFERENCE
  object s last assigned on line 4 could be null and is dereferenced at line 5
```

Beware of the inevitable false positives!



Analysis Dashboards

A holistic approach to quality: SonarQube

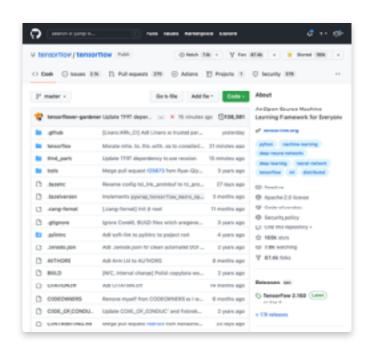








Let's look at a real project using SonarQube: TensorFlow





What analysis tools should I use?



The best QA strategies employ a combination of tools

How Many of All Bugs Do We Find? A Study of Static Bug Detectors

Andrew Habib andrewahabib@gamil.com Department of Computer Science TU Darristad Germany

ABSTRACT

Static bug detectors are becoming increasingly popular and are widely used by professional software developers. While most work on bur detectors focuses on whether they find bucs at all, and on how many take positives they report in addition to legitimate warnings, the inverse question is often neglected: How many of all real-world bugs do storie bug detectors find? This paper addresses this question by studying the results of applying three widely used static bug detectors to an extended version of the Defects @ fixtaget that consists of 15 Java projects with 194 known bugs. To decide which of these bugs the took detect, we use a novel methodology that combines an automatic analysis of warnings and bugs with a manual validation of each condicate of a detected bug. The results of the study show that: (i) static bug detectors find a non-negligible amount of all bugs, (i.) different tools are mostly complementary to each other, and (iii) current bug detectors miss the large majority of the studied burs. A detailed analysis of burs missed by the scatic detectors shows that some bugs could have been found by variants. of the existing detectors, while others are domain-specific problems that do not match any existing bug pattern. These findings helppotential opers of such tools to assess their at lity, motivate and outline directions for future work or, statte bug detection, and provide a basis for future comparisons of static bug detection with other bug finding techniques, each as manual and nationated testing.

Michael Pradel michael@binocrvarianade Department of Computer Science TU Damestadt Germany

International Conference on Automated Software Engineering (ASE 118), Septrepher L-2, MIR, Monipolites France, ACM, New York, NY, USA, 11 pages, Importable Argunt LLCV-230101-230619.

1 INTRODUCTION

Finding refevent bugs it an important but difficult task. For average industry code, the number of bugs per 1,940 lines of rode has been estimated to range between 0.5 and 25 [31]. Even after years of deployment, software still certains unmotiond bugs. For example, studies wither Linux for, set slow that the average bug crassion in the learned for a surprisingly long period of 3.6 to 1.8 years [4, 24]. Unfortunately, a single bug can cause across haven, even if it has been subsisting for a long time without doing so, as evidenced by examples of software bugs that have caused linge eventual loss and over all the people [11, 25, 40].

Given the importance of finding software bugs, developers only on several approaches to reveal programming mistakes. One approach is to identify hugs during the development process, e.g., through pair programming or rade review. Another direction is testing, sugging from purely manual testing over semi-automated testing, e.g., was manually written but automatically excepted unit tests, to fairly automated testing, e.g., with UI-level testing tools. Once the software is deployed, runtime monitoring concreved so far minor down on, e.g., collect information about abnormal runtime

Tool	Bugs
Error Prone	8
Infer	5
SpotBugs	18
Total:	31
Total of 27 uni	que bugs

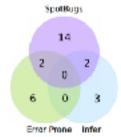


Figure 4: Total number of bugs found by all three static checkers and their overlap.

Summary

- Linters are cheap, fast, but imprecise analysis tools
 - Can be used for purposes other than bug detection (e.g., style)
- Conservative analyzers can demonstrate the absence of particular defects
 - At the cost of false positives due to necessary approximations
 - Inevitable trade-off between false positives and false negatives
- The best QA strategy involves multiple analysis and testing techniques
 - The exact set of tools and techniques depends on context

